

# CSCE 451/851

## Operating Systems Principles

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### Processes

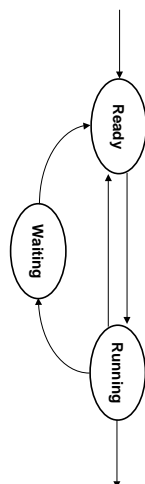
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<http://www.cse.unl.edu/~goddard/Courses/CSCE451>

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### Processes

- ◆ The basic agent of work, the basic building block
- ◆ Process characterization
  - Program code
  - Processor/Memory state
  - Execution state
- ◆ The state transition diagram



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## Process Actions

- ◆ Create and Delete
- ◆ Suspend and Resume
- ◆ Process synchronization
- ◆ Process communication

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## Physical v. Logical Concurrency

### Why is logical concurrency useful?

- ◆ Structuring of computation
- ◆ Performance

```
process P          system call Read()
begin
:                  begin
:                  StartIO(input device)
:                  WaitIO(interrupt)
:                  EndIO(input device)
:                  :
end P              end Read
```

» Single process I/O

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## Physical v. Logical Concurrency

### Performance considerations

#### ◆ Multithreaded I/O

```
process P
begin
:
  StartRead()
  <compute>
  Read(var)
:
end P

system call StartRead()
begin
  RequestIO(input device)
end StartRead

system call Read()
begin
  SignalReader(input device)
end Read
```

```
system process Read()
begin
  loop
    WaitForRequest()
    System_Read(var)
    WaitForRequestor()
  :
  end loop
end Read
```

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## Process Creation Paradigms

#### ◆ COBEGIN/COEND

```
cobegin
S1 ||
S2 ||
:
Sn
coend
```

#### ◆ FORK/JOIN

```
begin
:
  fork(foo)
:
  join(foo)
:
end

procedure foo()
begin
:
end foo
```

#### ◆ Explicit process creation

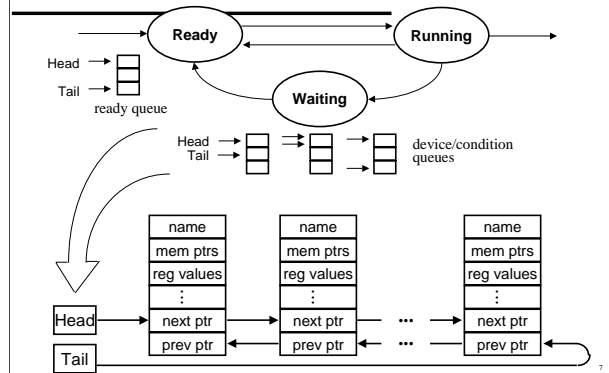
```
begin
:
P
:
end

process P
begin
:
end P
```

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## Process Scheduling

### Implementing and managing state transitions



## Why Schedule?

### Scheduling goals

- ◆ Example: two processes execute concurrently

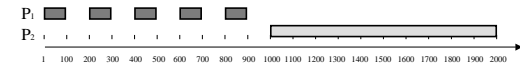
```

process P1
begin
  for i := 1 to 5 do
    <read a char>
    <process a char>
  end for
end P1

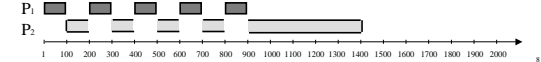
process P2
begin
  <execute for 1 sec>
end P2

```

- ◆ Performance without scheduling



- ◆ Performance with scheduling



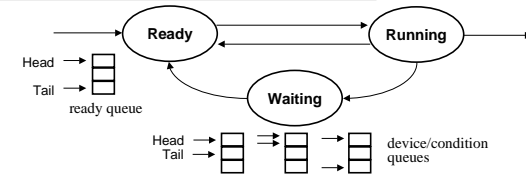
## Types of Schedulers

- ◆ Long term schedulers
  - » adjust the level of multiprogramming through admission control
- ◆ Medium term schedulers
  - » adjust the level of multiprogramming by suspending processes
- ◆ Short term schedulers
  - » determine which process should execute next

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## Short Term Scheduling

### When to schedule



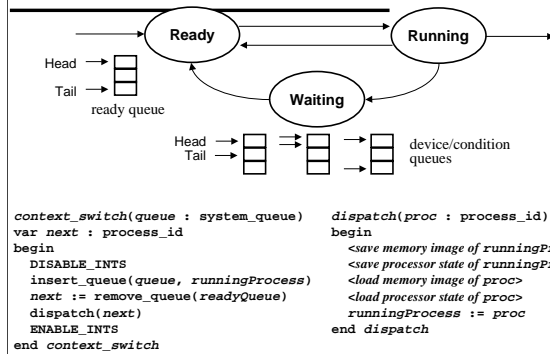
When a process makes a transition...

1. from *running* to *waiting*
2. from *running* to *ready*
3. from *waiting* to *ready*
- (3a. a process is *created*)
4. from *running* to *terminated*

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## Short Term Scheduling

How to schedule — Implementing a context switch

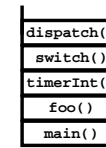


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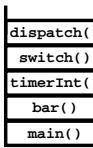
## Implementing a Context Switch

Dispatching

◆ Case 1: Preemption

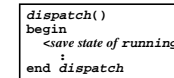


"running"

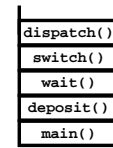


"next"

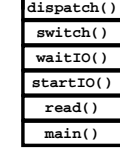
"running's" dispatch:



◆ Case 2: Yield

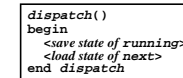


"running"



"next"

"next's" dispatch:



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## Producer/Consumer Implementation

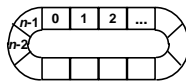
```

process Producer
  var c : char
begin
  loop
    <produce a character "c">
    while nextIn+1 mod n = nextOut do
      NOOP
    end while
    buf[nextIn] := c
    nextIn := nextIn+1 mod n
  end loop
end Producer

process Consumer
  var c : char
begin
  loop
    while nextIn = nextOut do
      NOOP
    end while
    c := buf[nextOut]
    nextOut := nextOut+1 mod n
    <consume a character "c">
  end loop
end Consumer

```

nextIn — nextOut



globals  
 buf : array [0..n-1] of char;  
 nextIn, nextOut : 0..n-1 := 0

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## Producer/Consumer Implementation with a shared counter

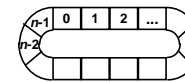
```

process Producer
  var c : char
begin
  loop
    <produce a character "c">
    while count = n do
      NOOP
    end while
    buf[nextIn] := c
    nextIn := nextIn+1 mod n
    count := count + 1
  end loop
end Producer

process Consumer
  var c : char
begin
  loop
    while count = 0 do
      NOOP
    end while
    c := buf[nextOut]
    nextOut := nextOut+1 mod n
    count := count - 1
    <consume a character "c">
  end loop
end Consumer

```

nextIn — nextOut



globals  
 buf : array [0..n-1] of char;  
 nextIn, nextOut : 0..n-1 := 0  
 count : integer := 0

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## Process Coordination

### Producer/Consumer systems

#### ◆ Example: Synchronous I/O

